

A Greedy Approach to Combinatorial Auctions

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Greedy allocation scheme

- ⌘ Assume single-minded bidders (one bid per bidder)
- ⌘ Order bids by descending average price per good
- ⌘ Allocate greedily going down the list
- ⌘ Each winning bid pays
(number of goods in bid) *
(average-price-per-good in highest displaced bid)
- ⌘ Losing bidders pay nothing

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Example

- ⌘ Bids: Red bids 10 from {a}, Green bids 19 for {a,b}, Blue bids 8 for {b}
- ⌘ GVA: Green gets {a,b}; pays 10 to Red, 8 to Blue
- ⌘ Greedy allocation: Red gets {a}, Blue gets {b}
- ⌘ GVA-style payment: Red pays 19-8=11, Green and Blue pay nothing. Note: Red's payoff negative.
- ⌘ New-style payment: Red pays 9.5, Green and Blue nothing

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Why study this greedy method?

- ⌘ A simple test case
- ⌘ $1/k$ lower-bound on the approximation ratio $1/k$, where k is number of goods
- ⌘ Can replace scoring function a/s by $a/?s$, in which case lower bound drops to $1/?k$
- ⌘ In practice usually better
- ⌘ Proof techniques very general and apply to other CAs
- ⌘ In particular, to an anytime version of the greedy method

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Sufficient conditions for incentive compatibility with single-minded bidders

1. **Exactness:** You get what you asked for or you get nothing
 2. **A-monotonicity:** If you win, you'd also win if you asked for less or offered more money or both
- Lemma:** Exactness and A-monotonicity imply the existence of V_c , above which you win and below which you lose
3. **Criticality:** If you win you pay V_c
 4. **P-monotonicity:** If you win, you'd pay less if you asked for less
 5. **Participation:** if you don't win you pay nothing
- Theorem:** Exactness, A-monotonicity, Criticality, P-monotonicity and Participation ensure incentive compatibility

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Further truths

- ⌘ The greedy scheme is incentive compatible for single-minded bidders
 - Proof: it satisfies the axioms
- ⌘ An "anytime" CA
 - Perform chronological backtracking to find better solutions (in the limit, the optimal one)
 - At each level can stop the algorithm and calculate new-style payments
 - Axioms still hold, hence does incentive compatibility
- ⌘ Beyond single-minded bidders
 - No payment scheme can make the greedy mechanism incentive compatible for complex bidders

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